

*Application*  
*for*  
*United States Patent*

*To all whom it may concern:*

*Be it known that I, Robert C. Klein, Jr., a citizen of the USA, have invented a new and useful*

*Self-Configuring Processing Element*

*of which the following is a full and clear description:*

**SELF-CONFIGURING PROCESSING ELEMENT**CLAIM OF PRIORITY

[0001] This application claims priority to, and incorporates by reference in its entirety, the U.S. provisional patent application no. 60/398,149, filed July 23, 2002.

FIELD OF THE INVENTION

[0002] The present invention relates generally to a configurable processing block and, more specifically, to a self-configuring processing element for providing arbitrarily wide application-specific instruction set extensions to a standard Instruction Set Architecture microcontroller in a semiconductor device.

BACKGROUND OF THE INVENTION

[0003] Various forms of configurable processing elements have been implemented in Field Programmable Gate Arrays (FPGAs) and Complex Programmable Logic Devices (CPLDs). In traditional FPGA and CPLD architectures, configurable processing elements include Look-Up Table (LUT)-based and/or multiplexer-controlled logic elements.

[0004] One problem with devices using conventional configurable processing elements is configuration latency. In such devices, every aspect of the device is programmed after the chip is powered on, including every logical function and every connection point for a given application. Each of these functions and connection points must be set by values contained in a configuration bit stream. As the size of the configuration bit stream increases, the delay in loading the configuration

bit stream increases. Since the configuration bit stream is typically loaded serially, the configuration latency is directly proportional to the size of the configuration file.

[0005] Another problem that results from an increase in the size of the configuration bit stream is that the cost of a solution using devices with conventional configuration processing elements increases. As the number of functions and connection points increases, larger configuration files are required. Larger configuration files require larger external memories in which to store the files. Thus, as the size of the configuration bit stream increases, the size and cost of the external memory storing the configuration bits increases as well.

[0006] Yet another problem with devices using conventional configurable processing elements is that the entire device must be configured, or reconfigured, in one process. Conventional configurable processing elements are not capable of performing either a partial reconfiguration or a pipelined reconfiguration in typical operation.

[0007] While devices using conventional configurable processing elements may be suitable for the particular purpose to which they were designed, they are not suitable for providing arbitrarily wide, application-specific instruction-set extensions to a standard Instruction Set Architecture (ISA) microcontroller.

### SUMMARY OF THE INVENTION

[0008] In view of the foregoing disadvantages inherent in the known types of configurable processing elements, the self-configuring processing element according to the present invention substantially departs from the conventional concepts and designs of the prior art. In so doing, the self-configuring processing element provides an apparatus developed to solve one or more of the problems described above. For example, a preferred embodiment of the self-configuring processing

element may provide arbitrarily wide, application-specific instruction set extensions to a standard ISA microcontroller in a semiconductor device.

**[0009]** The general purpose of the present invention, which will be described subsequently in greater detail, is to provide a new self-configuring processing element that has many of the advantages of conventional configurable processing elements and novel features that result in a new self-configuring processing element.

**[0010]** In a preferred embodiment of the present invention, a processing element includes a system bus interface, an instruction handler, an input router and conditioner electrically connected to the system bus interface and the instruction handler, an ALU electrically connected to the input router and conditioner, a memory electrically connected to the input router and conditioner, and an output router electrically connected to the ALU, the memory and the input router and conditioner.

**[0011]** In an embodiment, the system bus interface and instruction handler include a connection to a system bus having a plurality of address lines and a plurality of data lines, an address decoder, connected to one or more of the plurality of address lines, for determining whether the processing element is selected by comparing a value contained on the one or more address lines with a decoding value and asserting an enable flag when the processing element is selected, an instruction register, connected to one or more of the plurality of address lines and one or more of the plurality of data lines, for storing the values contained on the one or more address lines and the one or more data lines when the enable flag is asserted, and a state machine, connected to the instruction register, for configuring the processing element based on at least one of the stored address value and the stored data value.

**[0012]** In an embodiment, the input router and conditioner include a first input path connected to an output of a first input processing element, a second input path connected to an output

of a second input processing element, a third input path connected to an output of a third input processing element, one or more multiplexers for determining a data value, an address/data value, and a carry bit, and circuitry for selectively performing one or more operations on at least one of the data value and the address/data value and the carry bit. In an embodiment, the input router and conditioner further includes a fourth input path connected to a feedback path and/or a system bus.

**[0013]** In an embodiment, the one or more operations include performing a bit shift operation on at least one of the data value and the address/data value, incrementing at least one of the data value and the address/data value, decrementing at least one of the data value and the address/data value, storing at least one of the data value and the address/data value, and passing through at least one of the data value and the address/data value.

**[0014]** The one or more multiplexers may include a first multiplexer for determining a first portion of the data value, a second multiplexer for determining a second portion of the data value, a third multiplexer for determining a first portion of the address/data value, a fourth multiplexer for determining a second portion of the address/data value, and a fifth multiplexer for determining the carry bit. The first portion of the data value and the second portion of the data value may be of equal width. The first portion of the address/data value and the second portion of the address/data value may be of equal width.

**[0015]** In an embodiment, the first input processing element is located along an x-axis with reference to the processing element, the second input processing element is located along a y-axis with reference to the processing element, and the third input processing element is located in a diagonal direction with reference to the processing element.

**[0016]** In an embodiment, the output routing block includes a first output path connected to an input of a first output processing element, a second output path connected to an input of a second

output processing element, and a third output path connected to an input of a third output processing element. The output router may further include a fourth output path connected to a feedback path and/or a data bus. In an embodiment, the first output processing element is located along an x-axis with reference to the processing element, the second output processing element is located along a y-axis with reference to the processing element, and the third output processing element is located in a diagonal direction with reference to the processing element.

**[0017]** In a preferred embodiment, a method of configuring a processing element includes providing an address value and a data value to the processing element, decoding the address value, determining from the decoded address value whether the processing element is selected, if the processing element is selected, storing at least a portion of the address value and the data value, loading the stored address value and the stored data value into a state machine associated with the processing element, and configuring, by the state machine, the processing element based on the stored address value and the stored data value. The configuring step may include enabling one or more components of the processing element, and determining the routing of one or more multiplexers within the processing element. The configuring step may further include storing one or more values, determined by at least one of the stored address value and the stored data value, in a memory.

**[0018]** In an alternate embodiment, a method of configuring a processing element includes providing an address value to the processing element, decoding the address value, determining from the decoded address value whether the processing element is selected, if the processing element is selected, storing at least a portion of the address value, loading the stored address value into a state machine, and configuring, by the state machine, the processing element based on the stored address value.

**[0019]** In an alternate embodiment, a processing element includes an input block and an output block. The input block includes a first input path connected to an output of a first input processing element, a second input path connected to an output of a second input processing element, a third input path connected to an output of a third input processing element. The output block includes a first output path connected to an input of a first output processing element, a second output path connected to an input of a second output processing element, and a third output path connected to an input of a third output processing element. In an embodiment, the input block further includes a fourth input path connected to a feedback path and/or a system bus. In an embodiment, the first input processing element is located along an x-axis with reference to the processing element, the second input processing element is located along a y-axis with reference to the processing element, and the third input processing element is located in a diagonal direction with reference to the processing element. In an embodiment, the output block further includes a fourth output path connected to a feedback path and/or a system bus. In an embodiment, the first output processing element is located along an x-axis with reference to the processing element, the second output processing element is located along a y-axis with reference to the processing element, and the third output processing element is located in a diagonal direction with reference to the processing element.

**[0020]** There has thus been outlined, rather broadly, the more important features of the invention in order that the detailed description thereof may be better understood, and in order that the present contribution to the art may be better appreciated. There are additional features of the invention that will be described hereinafter.

**[0021]** In this respect, before explaining at least one embodiment of the present invention in detail, it is to be understood that the invention is not limited in its application to the details of

construction and to the arrangements of the components set forth in the following description or illustrated in the drawings. The invention is capable of other embodiments and of being practiced and carried out in various ways. Also, it is to be understood that the terminology used herein is for the purpose of the description and should not be regarded as limiting.

#### BRIEF DESCRIPTION OF THE DRAWING

[0022] Various other objects, features and attendant advantages of the present invention will become fully appreciated as the same becomes better understood when considered in conjunction with the accompanying drawings, in which like reference numbers designate the same or similar parts throughout the following text.

[0023] FIG. 1 depicts an exemplary embodiment of a self-configuring processing element according to an embodiment of the present invention.

[0024] FIG. 2 is a flowchart illustrating exemplary steps in a method of configuring the processing element.

[0025] FIG. 3 depicts an exemplary use of a group of self-configuring processing elements in a two-dimensional toroidal interconnect structure.

#### DETAILED DESCRIPTION OF THE INVENTION

[0026] Before the present methods are described, it is to be understood that this invention is not limited to the particular methodologies or protocols described, as these may vary. It is also to be understood that the terminology used in the description is for the purpose of describing the particular versions or embodiments only, and is not intended to limit the scope of the present invention which will be limited only by the appended claims. In particular, although the present invention is



described in conjunction with a silicon-based electrical circuit, it will be appreciated that the present invention may find use in any electrical circuit design.

**[0027]** It must also be noted that as used herein and in the appended claims, the singular forms “a”, “an”, and “the” include plural references unless the context clearly dictates otherwise. Thus, for example, reference to a “processing element” is a reference to one or more processing elements and equivalents thereof known to those skilled in the art, and so forth. Unless defined otherwise, all technical and scientific terms used herein have the same meanings as commonly understood by one of ordinary skill in the art. Although any methods similar or equivalent to those described herein can be used in the practice or testing of embodiments of the present invention, the preferred methods are now described. All publications mentioned herein are incorporated by reference. Nothing herein is to be construed as an admission that the invention is not entitled to antedate such disclosure by virtue of prior invention.

**[0028]** Turning now descriptively to the drawings, FIG. 1 illustrates a self-configuring processing element **100**, which may include the System Bus Interface and Instruction Handling (SBI) block **110**, the Input Routing and Conditioning (IRC) block **120**, the Arithmetic Logic Unit (ALU) block **130**, the Memory block **140**, and/or the Output Routing block **150**.

**[0029]** The SBI block **110** accepts address, data, and control information from one or more microcontrollers, microprocessors, digital signal processors and/or state machines via a system bus **114**. The one or more microcontrollers, microprocessors, digital signal processors, and/or state machines may reside in the same electrical circuit as the processing element **100**, or it may be external to the electrical circuit. Although FIG. 1 illustrates a 32-bit system bus, system busses of other sizes may be used. The SBI block **110** may include a cell ID address decoder **111**, a register for holding appropriate bits from the system address bus **115** and system data bus **116**, a state machine

for sequencing through processing element initialization and instruction set-up tasks, and/or tri-state buffers **113** for controlling data flow to and from the system bus **114** and/or for feedback within the processing element **100**. The above-described register and state machine are collectively represented by block **112** in FIG. 1.

[0030] A specific range of binary addresses may be assigned to each processing element integrated into a system. The cell ID address decoder **111** of the SBI block **110** may respond to a specific range of addresses in the address field of the system bus **114** that are defined for the particular instance in which the cell ID address decoder **111** is located. If the information present on the system bus **114** falls within the range, the cell ID address decoder **111** may enable the Instruction Register, Decode, and State Machine logic block **112** via an enable signal. The Instruction Register, Decode, and State Machine logic block **112** may respond by decoding the information from the address bus **115** and the data bus **116** in order to perform one or more of several actions. These actions may include, but are not limited to, the following:

1.       WRITEMEM: This function may write data from the data bus **116** to a given location in the Memory block **140**. The address of the location to be modified may be determined by information from the address bus **115**. This command may be used to create a full-custom instruction by specifying the contents of the Memory block **140** for Look-Up Table (LUT) logical functions.
2.       READMEM: This function may drive the contents of the Memory block **140** onto the system bus. The address of the location to be read may be determined by information from the address bus **115**.
3.       READALU: This function may drive the contents of the ALU block **130** onto the data bus **116**.

4. READBUS: This function may drive a copy of one of the input busses **121** or output busses **152** onto the data bus **116**. The source bus (i.e., whether an input **121** or output bus **152** is read) may be determined by information from the address bus **115**.

5. WRITEBUS: This function may drive one of the input busses **121** or output busses **152** with the data on the data bus **116**. The destination bus may be determined by information from the address bus **115** which may drive the select lines of the Output Multiplexers **151**.

6. WRITEINST: This function may initialize the state machine **112** in the SBI block **110**. The addressed processing element **100** may perform a series of actions controlled by the state machine **112** that result in the processing element **100** being configured to perform one of a predetermined set of instructions. Information on the address bus **115** may determine which instruction is used to configure the processing element **100**. The predetermined set of instructions may be further refined by the contents of the data bus **116**. For example, a command may be issued to instruct the processing element **100** to create a “Multiply by \$7E” instruction (a hexadecimal multiply-by-a-constant function). The selection of the “multiply-by-a-constant” configuration may be encoded in the address bus **115**, while the “\$7E” (i.e., the specific constant to multiply by) may be read from the data bus **116**.

7. SELECTIN: This function may determine one or more sources for subsequent input data **124-127** and carry-in **128** signals for the processing element **100**. The one or more sources may be determined by information in the address or data fields of the system bus **114**. The routing may be performed by the Input Multiplexers **123**.

8.       SELECTOUT: This function may determine one or more destinations for subsequent output data **152** and **153** and the carry-out signal **132** for the processing element **100**. The one or more destinations may be determined by information in the address or data fields of the system bus **114**.

9.       SELECTMEM: This function may configure the processing element **100** and its associated Memory block **140** to be one of a pre-determined set of memory functions. These memory functions may include, but are not limited to, Static Random Access Memory (SRAM), First-In-First-Out (FIFO), Last-In-First-Out (LIFO), Content Addressable Memory (CAM), or a shift register. The selection of the function for the Memory block **140** may be made based on information in the address or data fields of the system bus **114**.

**[0031]** The SBI block **110** is not limited to the construction set forth above. Variations on this block may include, but are not limited to, alternate system bus interface architectures resulting from different system busses being used, including a system bus where information is passed over shared connections such as the Toroidal Input Busses **121**, alternate methods of decoding and using the information from the data bus **116**, the address bus **115** and control signals, different bus word widths and data word widths, and support for modified or different instructions by the state machine **112**. The microcontrollers, microprocessors, digital signal processors and/or state machines controlling the system bus may be either on-chip or off-chip. The instructions and data may also be supplied by other processing elements connected, either directly or indirectly, to the self-configuring processing element **100**.

**[0032]** FIG. 2 is a flowchart illustrating exemplary steps in a method of configuring the processing element **100**. First, an address value and/or a data value may be provided **200** to the processing element **100**. The address value may be decoded **205**, and a determination may be

made **210** from the decoded address value as to whether the processing element is selected. If the processing element **100** is selected, at least a portion of the address value and/or the data value may be stored **215**. The stored address value and/or the stored data value may be loaded **220** into a state machine associated with the processing element **100**. The state machine may configure **225** the processing element **100** based on the stored address value and/or the stored data value. This configuration may include, but is not limited to, setting enable flags and multiplexer selects, defining memory locations in the Memory block **140**, and determining the function to perform in the ALU **130**.

[0033] Returning to FIG. 1, the Input Routing and Conditioning block **120** may select and connect the available inputs to the ALU block **130** and the Memory block **140** via Input Multiplexers **123**. In addition, the IRC block **120** may include circuitry for registering, shifting, incrementing, and/or decrementing the inputs received or loaded. Such circuitry is collectively represented by block **122** of FIG. 1. The configuration of the Input Multiplexers **123** and the specific action to be performed on the incoming data may be determined by information in the Instruction Register, Decode and State Machine logic block **112** in the SBI block **110**.

[0034] A method of processing an exemplary instruction will now be described in order to show the operation of the IRC block **120**. The SBI block **110** may receive information from the address bus **115** requesting that the processing element **100** implement a “multiply by a constant” function. The State Machine **112** in the SBI block **110** may load the constant to be multiplied from the data bus **116** into a register in the circuitry of block **122** that has an output sent to one input to the ALU block **130**. The ALU **130** may be set to accumulation mode (add-to-output) by the SBI block **110**. The incrementor in the circuitry of block **122** may then, starting from zero, supply address information to the memory, which may be SRAM or other appropriate memory, in the

Memory block **140**. The State Machine **112** in the SBI block **110** may then cycle through one state for each location in the Memory block **140**. In a preferred embodiment, 256 memory locations are used, and the State Machine **112** may cycle through 256 states. In each state, the value stored in the register in the IRC block **120** may be added to the output of the ALU **130**, the counter in the circuitry of block **122**, which is connected to the address inputs of the Memory **140**, may increment, and the selected location in Memory **140** may be written with the accumulated data from the output of the ALU **130**. When this process is completed and the instruction is executed, the Memory **140** may respond by outputting a result equal to the constant multiplied by a value on the address lines of the Memory **140**.

[0035] In a preferred embodiment, this function may be initialized by a single command received from the system bus **114**. Once the command is issued, the initialization procedure may proceed without the intervention or control of the system bus **114** or any external device. The lack of the need for direct control over the initialization procedure may allow the system bus **114** to be used to perform other tasks instead of monitoring particular processing elements or waiting for the initialization procedure to complete. In this manner, the configuration latency inherent in devices using conventional configurable processing elements may be reduced in devices incorporating the present invention. Of course, systems using control by the system bus **114**, although not required, may be included in the scope of the present invention.

[0036] The connections between the IRC block **120** and the ALU/Memory block **130** will now be described. In a preferred embodiment, as shown in FIG. 1, there may be, for example, four separate busses that are used to form the data and address inputs to the Memory **140**. Each bus may also be used to form the X and Y inputs of the ALU **130**. Each bus, in a preferred embodiment, may be four bits wide. Alternate widths may be selected for each bus individually without limitation. In

addition, a carry-in signal may be passed to the ALU **130**. The carry-in signal may also be used as the input to the least significant bit of the shifter/counter circuitry **122** in the IRC block **120**. The shift out signal of the most significant bit of the shifter/counter circuitry **122** may be an additional single-bit output that is presented to the Output Routing block **150** for direction to its ultimate destination (if any).

[0037] Variations on these signals may include altering the width of the input busses **121** and/or selection circuitry **122**, changing the method of encoding, decoding and routing the input busses **121** to the outputs of the circuitry **122**, and modifying the logical structure of the internal shifter/counter circuitry **122**. Each of these modifications will be apparent to one of skill in the art and are considered to be within the scope of this invention.

[0038] The ALU block **130** may receive inputs **124-127** from the IRC block **120** and perform operations on such inputs **124-127** based on the information in the Instruction Register, Decode and State Machine logic **112** in the SBI block **110**. The ALU block **130** may include an eight-bit ALU (with 16 outputs to account for overflow and accumulation). The IRC block **120** may determine the sources for the various inputs **124-127** to the ALU **130**. Variations on the ALU block **130** may include, without limitation, ALUs of different widths, different input bus widths, variations in the functions performed by the ALU, and/or the potential sources and destinations of data operated on by the ALU. Each of these modifications, including designing ALUs and the functions performed by ALUs, will be apparent to one of skill in the art and are considered to be within the scope of this invention.

[0039] The Memory block may receive inputs **124-127** from the IRC block **120** and perform operations on such inputs **124-127** based on the information in the Instruction Register, Decode and State Machine logic **112** in the SBI block **110**. The Memory block **140** may include a

memory. In a preferred embodiment, the Memory block **140** may include a dual-port 256x8 SRAM cell (with separate read and write data ports, but a common address port). Additional logic in the IRC block **120** may be used to make the memory element operate as, for example, a FIFO, LIFO, CAM, or LUT. In the LUT mode, any logical function of eight inputs may be realized in the memory element. After a desired function is loaded into the memory, as determined by a microcontroller and received by the SBI block **110** via a system bus, the data for performing the function may be supplied by the IRC block **120** to the memory. Based on the information stored in the memory, any logical function may be performed. Alternate memories including, without limitation, DRAMs, FLASH, and EEPROMs may be used instead of SRAM. In addition, the memory may be of different size and may have a different read/write port configuration.

[0040] The Output Routing block **150** may receive data from the outputs of the ALU block **130** and the Memory block **140** and route the data to one or more of a plurality of destinations. The specific destinations to be selected may be determined by information in the Instruction Register, Decode and State Machine logic **112** in the SBI block **110**. In a preferred embodiment, the Output Routing block **150** may include, for example, four byte-wide (eight-bit) four-to-one multiplexers **151** that select sources for three output busses **152** and one feedback bus **153**. A separate two-to-one multiplexer **151** may be provided to determine whether the most significant bit **129** of the shifter/counter circuitry **122** of the IRC block **120** or the carry out bit **132** from the ALU block **130** is used as a source for the three output busses **152** and the feedback bus **153**. The SBI block **110** may select the source passed through each multiplexer **151** based on the decoded instruction received from the system bus **114**. Details of the connections to and from the Output Routing block **150** will be set forth later in this document.



[0041] Variations in the Output Routing block **150** may include changes to the quantity and word widths of the inputs and outputs **152** and **153**, the decoding of the potential sources and destinations **152** and **153**, or the granularity of control (i.e., the number of bits that may be selected from each source and combined and sent to a given destination). Each of these modifications will be apparent to one of skill in the art and are considered to be within the scope of this invention.

[0042] In a preferred embodiment, a number of different types of connections may be present with respect to a processing element **100**. These connections may include connections via the system bus **114** to other system resources, such as one or more microcontrollers, microprocessors, digital signal processors, state machines, input/output pins, communication ports, and/or bulk memory blocks, connections from one processing element **100** to other processing elements, and connections within an individual self-configuring processing element **100**.

[0043] Referring to FIG. 1, the system bus **114** may allow information and data to be sent to and from the self-configuring processing element **100**. The system bus **114** may be connected to on-chip and/or external functional blocks including, without limitation, one or more microcontrollers, microprocessors, digital signal processors, state machines, input/output pins, communication ports, and/or memory blocks. The system bus **114** may enable data, control, configuration and status information to be passed into and out of a logic fabric created by an array of processing elements, such as that illustrated in FIG. 3. The system bus **114** may be any microprocessor bus architecture used by those skilled in the art. Such busses are commonplace in CPUs, embedded microcontrollers, digital signal processors, and most application-specific integrated circuits (ASICs). The system bus **114** may contain address, data and control signals. The address signals may be used to determine the devices and/or locations on the system bus **114** that have been selected to transmit or receive data in a given system cycle. Data signals may be used to transfer information over the system bus **114**.

Control lines may include such signals as read/write, clock, reset, and enables that may be used for supervisory and/or timing purposes.

**[0044]** The many potential sources and destinations for the signals on the system bus **114** may require long, physically robust connections and additional buffering and/or drivers for the most heavily loaded signals. Since all logical and electrical functional blocks attached to the system bus **114** share these connections, a supervising program, processor or state machine may be used to determine which blocks send and receive data and in which order. To this end, a supervising program, processor or state machine may arbitrate simultaneous requests for the use of resources in order to avoid conflicts or bus contention.

**[0045]** In a preferred embodiment, the system bus **114** uses the ARM Microprocessor Bus Architecture (AMBA) as specified in the ARM AMBA manual (Doc No.: ARM IHI-0011, Issued: May 1999 by ARM Holdings plc, 90 Fulbourn Road, Cambridge CB1 9NJ, UK). This document describes an AHB (Advanced High-Performance Bus) and an APB (Advanced Peripheral Bus) that together comprise the system bus **114**. Only the APB attaches directly to a processing element **100**. A unique APB is used for each column of processing elements in a device. The columnar APB is addressed and activated by address information sent over the AHB. Information, such as configuration data and status information, and data may be passed between a microcontroller and the processing elements through this bus structure. The separation of control, implemented in the system bus **114**, and datapath, implemented in the interconnection of processing elements, permits a more efficient use of resources within devices incorporating one or more processing elements **100** according to the present invention.

**[0046]** In a preferred embodiment, each self-configuring processing element **100** may be connected to the system bus **114** through a columnar APB. All processing elements within a column

may share the address, data and control signals of the APB **114** associated with that column. The address signals of the APB **114** may be used to select one or more processing elements as the source or destination for the information carried in the data and control signals of the APB. In addition, the address lines may determine which data, configuration bits or memory locations within the one or more processing elements **100** are accessed.

[0047] Each individual columnar APB may be selectively connected to the AHB by decoding the address signals of the AHB. The columnar APBs may also serve as the connections to other system resources such as bulk memory blocks, input/output pins, and serial communication modules. Any configuration information needed by these other resources may also be sent and read-back across the columnar APBs.

[0048] With respect to the connections between processing elements, the preferred interconnection structure may be toroidal in nature, as described in a co-pending U.S. patent application entitled "Improved Interconnect Structure for Electrical Devices," filed July 23, 2003 with serial no. (not yet assigned), which is incorporated herein by reference in its entirety. The toroidal interconnect structure **300** may include, for example, three potential datapath sources **121** and, for example, three potential destinations **152** for each processing element **100**. These sources and destinations may include other processing elements **100**. Additional sources and destinations may include the system bus **114** and a feedback path **153** within a processing element **100**.

[0049] As shown in FIG. 3, the toroidal interconnect structure **300** may have x-direction (referred to herein as "horizontal" or "row") datapaths **310** and y-direction (referred to herein as "vertical" or "column") datapaths **320**. In addition, the toroidal interconnect structure **300** may have a diagonal, or effective "top left toward bottom right," datapath **330** that is also toroidal in nature. Other potential structural and functional variations may include providing a similar toroidal

interconnect along other diagonal paths, skipping multiple rows/columns, or simply creating the toroidal interconnect in fewer directions than is described herein (for example, a column-based, "vertical-only" toroidal interconnect.) Note that rows and/or columns are not necessarily skipped at edge elements, as an edge element may loop back to its nearest neighbor.

**[0050]** In FIG. 3, the terms "physical row" and "physical column" refer to the placement of a row or column, respectively, in a two-dimensional device layout. For example, the first physical row may be the row of processing elements **100** that are physically located at the top of the physical media. Sequentially subsequent physical rows may be adjacent to and below preceding physical rows. Likewise, physical columns may be arranged from left to right, where the first physical column is the leftmost column in the physical device. Other embodiments and orientations are possible within the scope of the invention.

**[0051]** In FIG. 3, the terms "row in toroid" and "column in toroid" refer to the placement of a row or column, respectively, in the three-dimensional representation embodied in a two-dimensional device layout. For example, the first row in the toroid may be the row of processing elements **100** physically located at the top of the physical media. A sequentially subsequent row in the toroid may be physically at least two rows below the preceding row in the toroid until an edge of the two-dimensional device is reached. At this point, sequentially subsequent rows in the toroid may be the "skipped" rows in the device ordered from the bottom of the device to the top. Likewise, columns in a toroid may be ordered by starting from the leftmost row, selecting every other row until the edge of the physical device is reached, and then selecting the "skipped" rows from right to left. Other embodiments and orientations are possible within the scope of the invention.

**[0052]** In the toroidal interconnect structure **300**, the potential inputs may be from a processing element along a y-axis (e.g., above), a processing element along an x-axis (e.g., to the

left), and a processing element diagonally disposed (e.g., above and to the left) from the processing element **100**. The data source for the processing element **100** may be selected from one or more of these potential source processing elements, the system bus **114**, or a feedback path **153**. The information from the selected data source **124-127** may be passed from the IRC block **120** into the ALU block **130** and the Memory block **140** via Input Multiplexers **123** and the shifter/counter circuitry **122** that may be controlled by the configuration of the processing element **100**.

**[0053]** The terms “above” and “to the left of” may not designate the physical two-dimensional relationships between processing elements. Instead, these terms may designate the placement of a processing element **100** within a three-dimensional toroidal interconnect structure **300**. In the physical device, the processing element **100** may be one or more rows or columns removed from the processing element which is “above” or “to the left of” the processing element **100**.

**[0054]** In a preferred embodiment incorporating the three-dimensional toroidal interconnect structure **300**, each processing element **100** may potentially output data to one or more of a processing element along a y-axis (e.g., below), a processing element along an x-axis (e.g., to the right), or a processing element diagonally disposed (e.g., below and to the right) from the processing element **100**. The output destinations may also include the system bus **114** or the feedback path **153** within the processing element **100**. The processing element **100** may drive one or more of these potential destinations **152** and **153** at the same time. The determination of which outputs **152** and **153** are driven by the Output Routing block **150** may be determined by the configuration of the processing element **100**.

**[0055]** The terms “below” and “to the right of” may not designate the physical two-dimensional relationships between processing elements. Instead, these terms may designate the

placement of a processing element **100** within a three-dimensional toroidal interconnect structure **300**. In the physical device, the processing element **100** may be one or more rows or columns removed from the processing element which is “below” or “to the right of” the processing element **100**.

[0056] With respect to the connections within a processing element **100**, the following connections represent an exemplary embodiment of the present invention. Variations may be made with regard to the connection paths including, without limitation, the width of the connection path, the source of the connection path, and the destination of the connection path. Each of these modifications will be apparent to one of skill in the art and are considered to be within the scope of this invention.

[0057] In a preferred embodiment, the system bus **114** may attach to the SBI block **110**. Address signals from the system bus **114** may be decoded by a cell ID address decoder **111** that may uniquely identify the address of the processing element **100**. In an embodiment, a number of address signals, for example, eight, may be attached from the system bus **114** to the IRC block **120**. These address signals **115** may be further grouped into sub-groups. In a preferred embodiment, each of two sub-groups may be four bits wide. These sub-groups may be individually selected by four-to-one Input Multiplexers **123** in the IRC block **120** that are controlled by the configuration contained in the SBI block **110** to determine the low-order (bits 3:0) and/or high-order (bits 7:4) inputs to the address inputs of the Memory **140** and/or the Y inputs of the ALU **130**. For example, the low-order address signals may be selected from a Toroidal Input Bus **121** and the high-order inputs may be selected from the system bus **114**.

[0058] In a preferred embodiment, if the processing element **100** recognizes its address on the system bus **114**, a number of data signals **116**, for example, eight, may be latched into the

Instruction Register, Decode and State Machine logic **112** in the SBI block **110**. The data signals **116** may also be passed to the IRC block **120**. The data signals **116** may be further grouped into sub-groups. In an embodiment, each of two sub-groups may be four bits wide. These sub-groups may be individually selected by four-to-one Input Multiplexers **123** in the IRC block **120** that are controlled by the configuration contained in the SBI block **110** to determine the low-order (bits 3:0) and/or high-order (bits 7:4) inputs to the data inputs of the memory and/or the X inputs of the ALU contained in the ALU/Memory block **130**. For example, the low-order input may be selected from the feedback path **153** and the high-order input may be selected from a toroidal input bus **121**.

**[0059]** In a preferred embodiment, the Output Routing block **150** may take the output from the Memory **140**, the output from the ALU **130**, and the output of the IRC block **120** as potential outputs to each of the processing element below (i.e., logically interconnected along a y-axis), the processing element to the right (i.e., logically interconnected along an x-axis) of and the processing element diagonally below and to the right of the processing element **100**, the system bus **114**, and the feedback path **153**. Optionally and preferably, the feedback path **153** is connected to the data path **116**. In a preferred embodiment, the output from the Memory **140** may be eight bits, the output from the ALU **130** may be sixteen bits, and the output of the IRC block **120** may be eight bits. These bit widths are exemplary only. Outputs of different size may be used within the scope of this invention. The selection of the bits to place on each output **152** and **153** may be performed via, for example, four eight-bit wide four-to-one Output Multiplexers **151** in the Output Routing block **150** and two banks of tri-state buffers **113** that are each eight bits in width (for the system bus **114** and feedback path **153** outputs). Preferably, a carry bit multiplexer **152** is also provided. The Output Multiplexers **152** preferably determine data value. The selection criteria may be decoded from the Instruction Register, Decode and State Machine logic **112** in the SBI block **110**. In addition, a ninth

bit may be sent to each of the three Toroidal Output Busses **152** and the feedback path **153** that contains either the carry-out **132** signal from the ALU **130** or the shift out signal **129** from the shifter/counter circuitry **122** in the IRC block **120**. The section criteria for the ninth bit may also be decoded from the Instruction Register, Decode and State Machine logic **112** in the SBI block **110**.

[0060] The Toroidal Input Busses **121** of a processing element **100** may, for example, be connected to the Toroidal Output Busses **152** of other processing elements. One method of connecting the processing elements is a toroidal interconnect structure **300** as shown in FIG. 3.

[0061] The connection paths internal to a processing element **100** described above represent only one method of interconnecting a self-configuring processing element **100**. Those skilled in the art will recognize that other methods of interconnecting the blocks of a processing element are evident based on this disclosure. Potential variations include changes to the number, connectivity and/or bus-widths of the processing element **100** to the Toroidal Input Busses **121**, the Toroidal Output Busses **152**, the feedback path signals **153**, and other internal busses. Changes to the bus widths may precipitate changes to the multiplexing structures of the IRC block **120** and the Output Routing block **150**. Changing the width and/or depth of the Memory **140** and the ALU **130** may also require changes to the fundamental architecture of the interconnection paths. Each of these modifications will be apparent to one of skill in the art and are collectively considered to be within the scope of the invention.

[0062] With respect to the above description, it is to be realized that the optimum dimensional relationships for the parts of the invention, including variations in size, materials, shape, form, function and manner of operation, assembly and use, are readily apparent to one of skill in the art, and all equivalent relationships to those illustrated in the drawings and described in the specification are intended to be encompassed by the present invention.



[0063] Therefore, the foregoing is considered as illustrative only of the principles of the invention. Further, since numerous modifications and changes will readily occur to those skilled in the art, it is not desired to limit the invention to the exact construction and operations shown and described, and accordingly, all suitable modifications and equivalents may be considered as falling within the scope of the present invention.